

How bad could it be?

The semilattice of definable sets in continuous logic

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CUNY Logic Workshop

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- Elementary extensions: $\mathbb{R} \oplus \mathbb{Q}^{\oplus \kappa}$, where $\mathbb{Q}^{\oplus \kappa}$ has $\{0, 1\}$ -metric.

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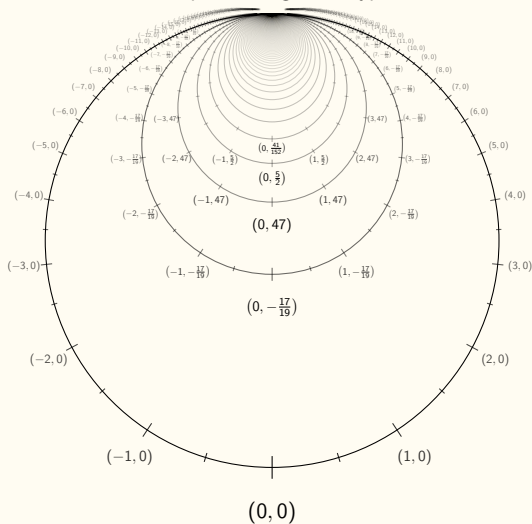
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- (H.) Any compact topometric space (X, τ, ρ) with open metric ρ is isomorphic to $S_1(T)$ for some strictly stable T .

Unique non-algebraic type



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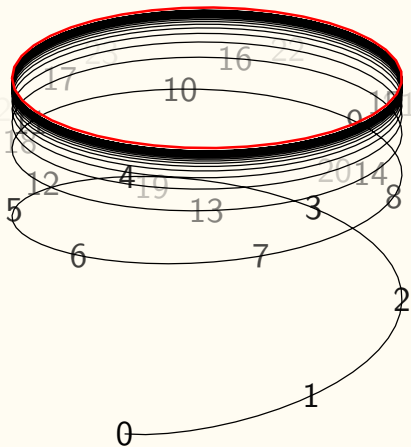
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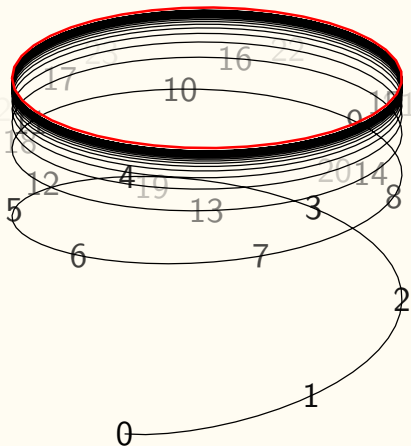
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- If T is ω -stable, then $S_n(A)$ always has a basis of definable neighborhoods. (T is *dictionary*.)

Many definable sets: $S_1(M)$, $M = (\mathbb{R}_{\geq 0}, \cos, \sin, d)$

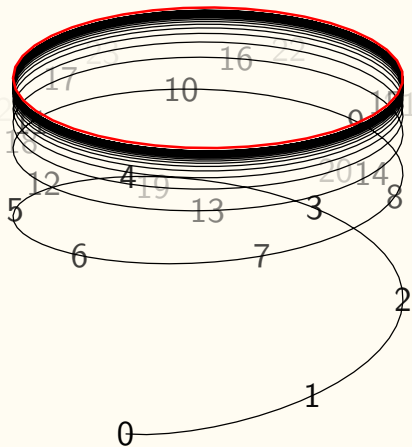


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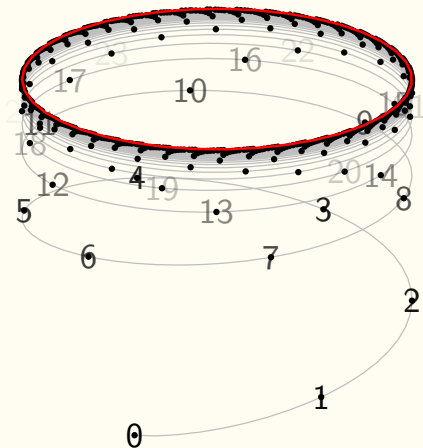
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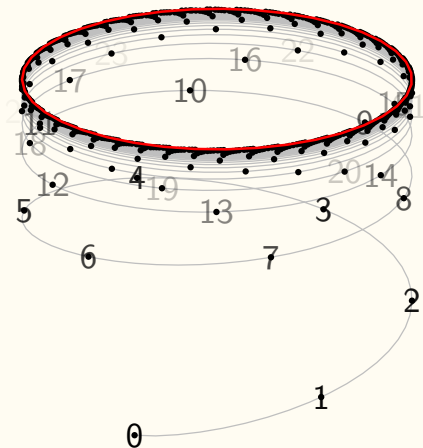


- $\text{Th}(M)$ is ω -stable, so has many definable sets (e.g. $\{x : \cos(x) \in F\}$ for any closed F).
- Metric on non-algebraic types is (roughly) path metric.

Few definable sets: $S_1(N)$, $N = (\mathbb{N}, \text{succ}, \cos, \sin, d)$

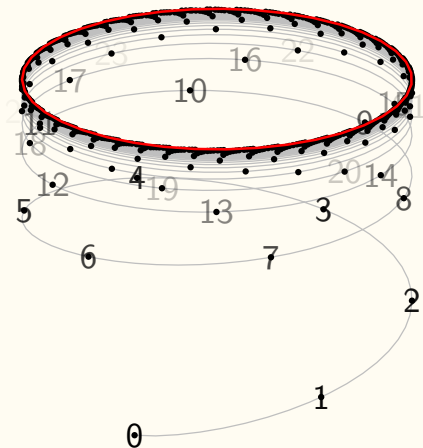


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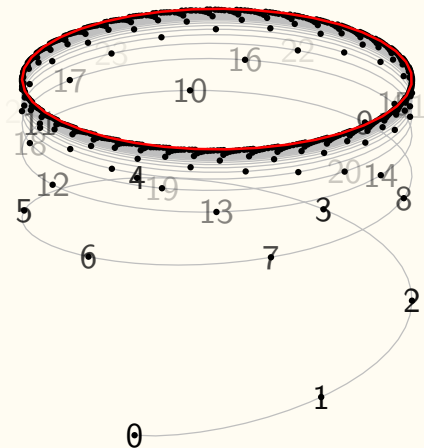
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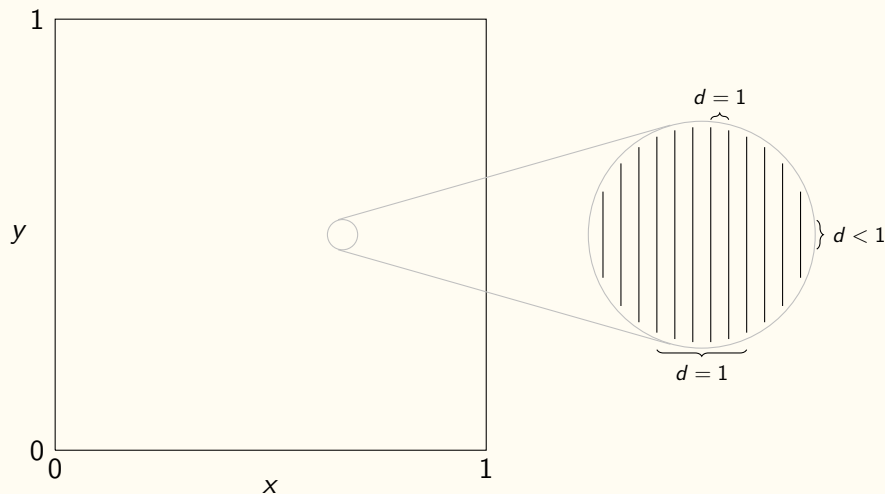
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- $\text{Th}(N)$ is superstable.
- Metric on non-algebraic types is discrete. Every definable set is either finite and algebraic or cofinite and co-algebraic.

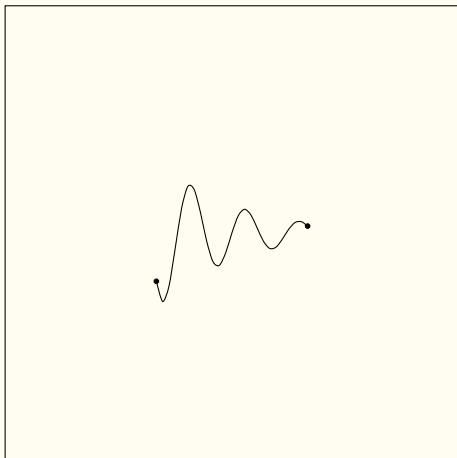
Many but not enough I



$$d((x_0, y_0), (x_1, y_1)) = 1 \text{ if } y_0 \neq y_1.$$

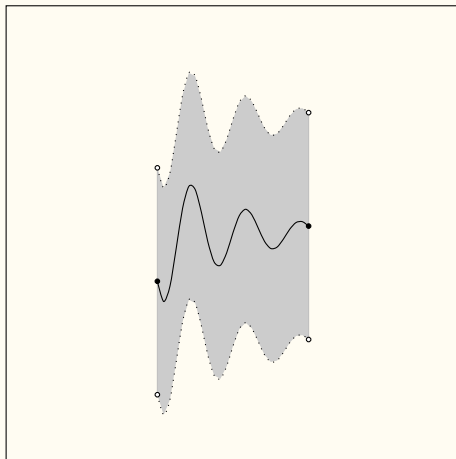
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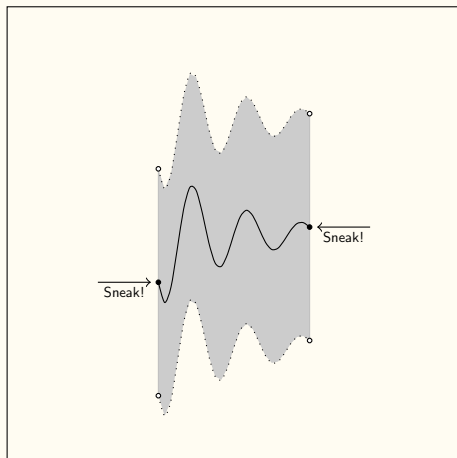
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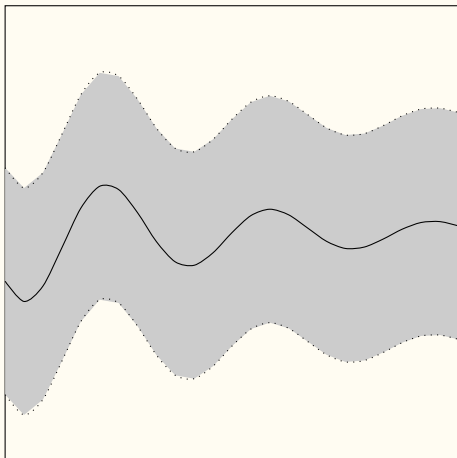
Closed F , with $F < \frac{1}{4}$.

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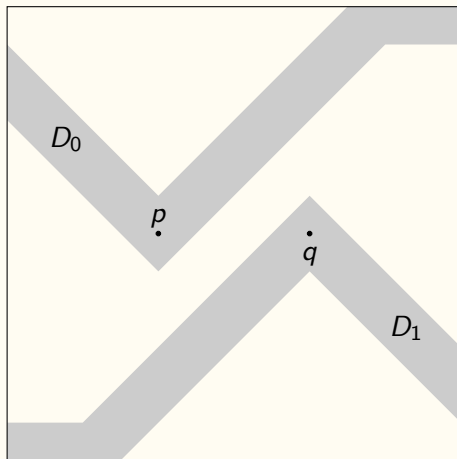
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Not definable.

Many but not enough III



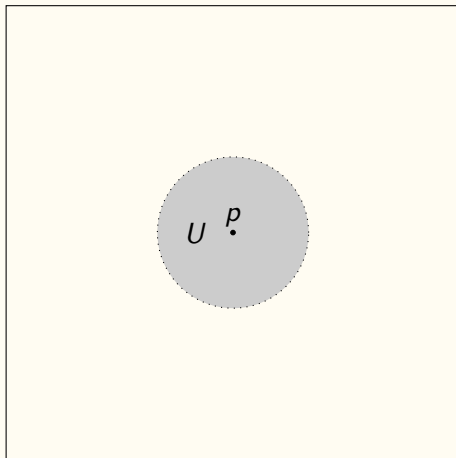
Definable set D , with $D < \frac{1}{4}$.

Many but not enough IV



Almost any two points are separated by disjoint definable neighborhoods.

Many but not enough V



There is no non-empty definable D with $D \subseteq U$.

The semilattice of definable sets

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Given a type space $S_n(T)$, the collection of definable subsets of it forms a bounded upper semilattice (\emptyset and $S_n(T)$ are always definable) under unions.

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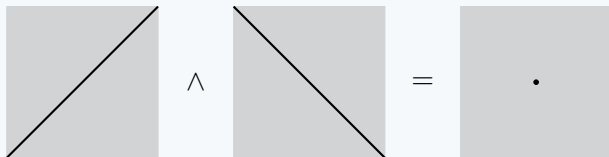
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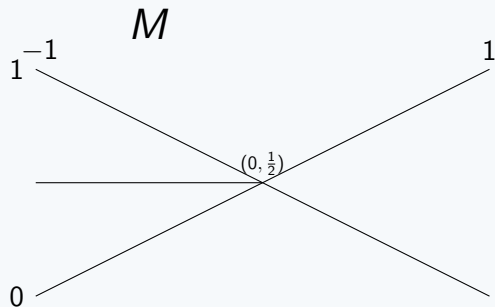
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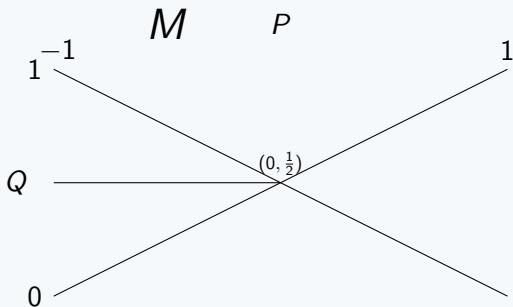


Prototypical example: Structure



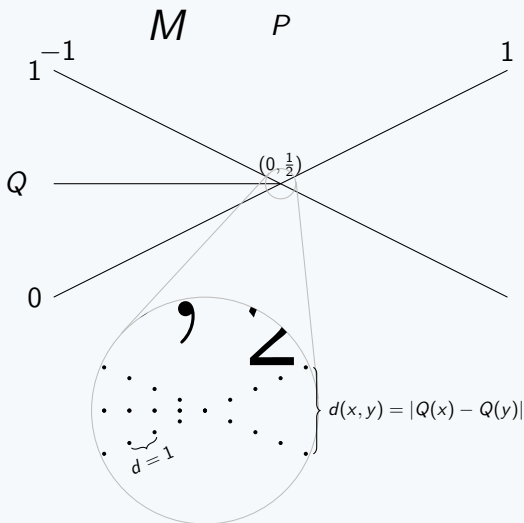
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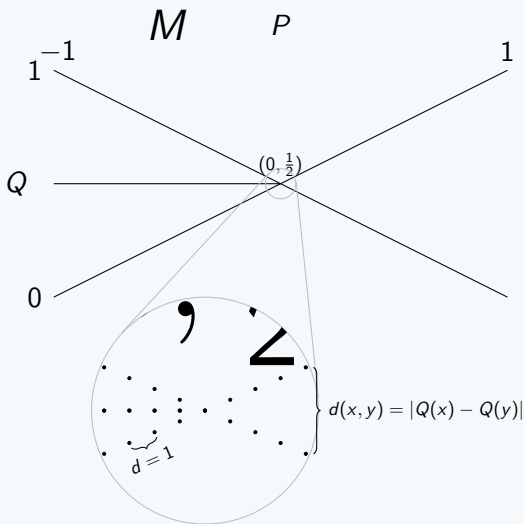
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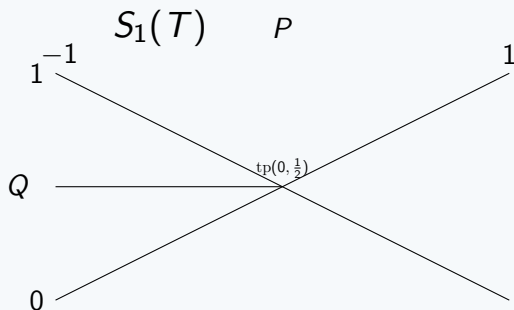
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- Let $T = \text{Th}(M)$.

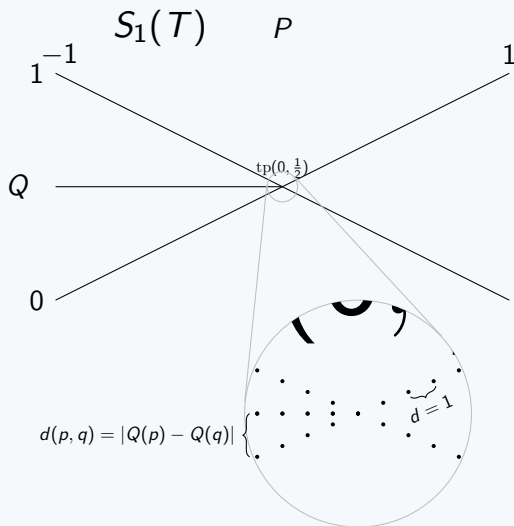
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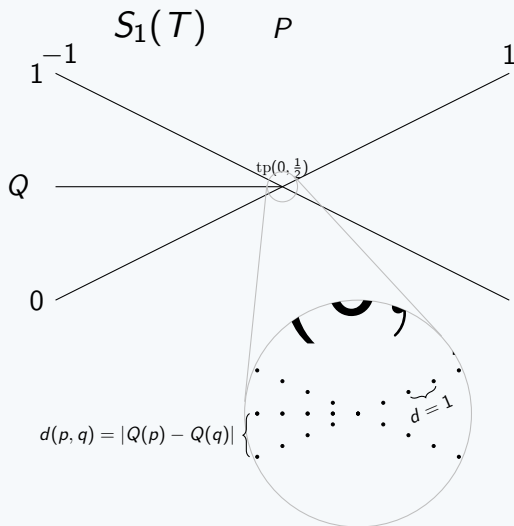
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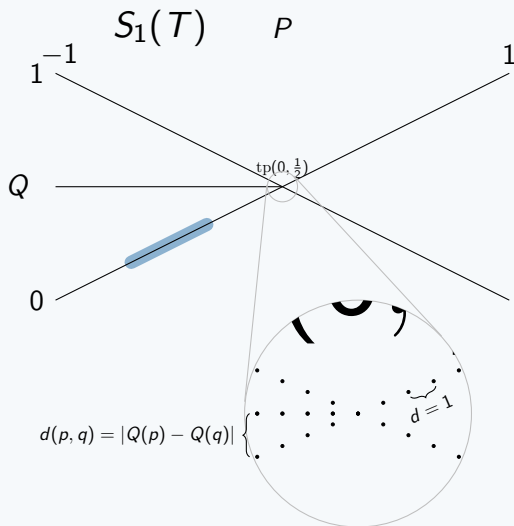
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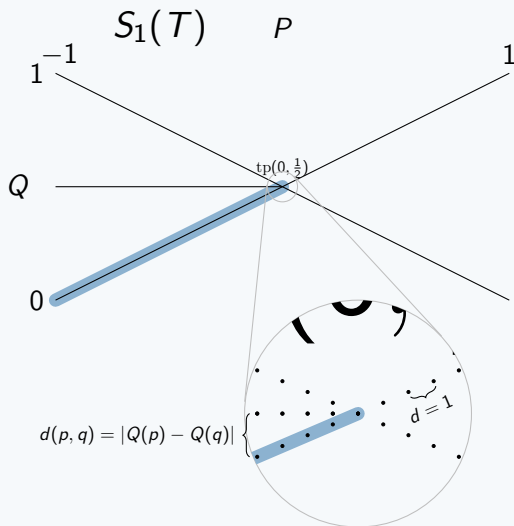
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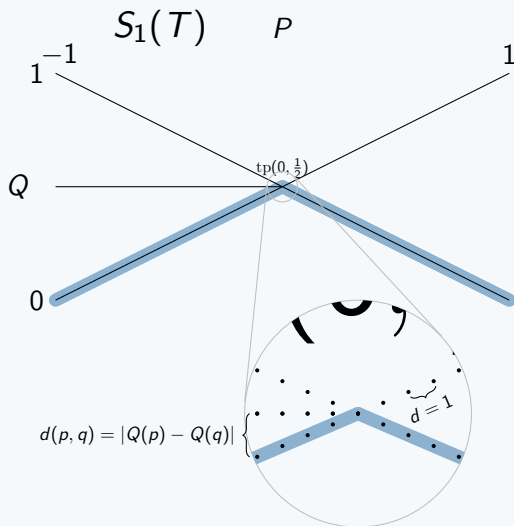
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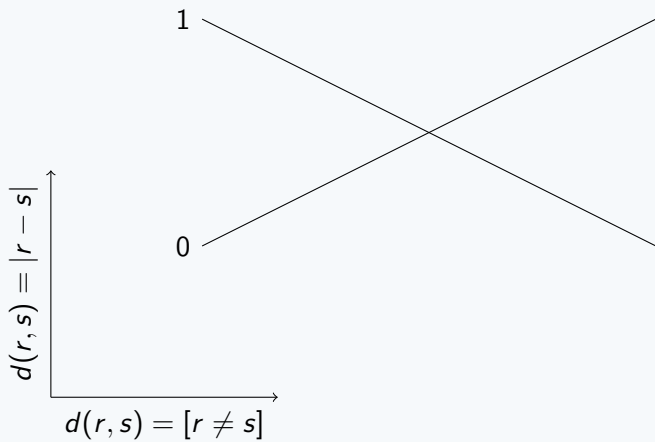
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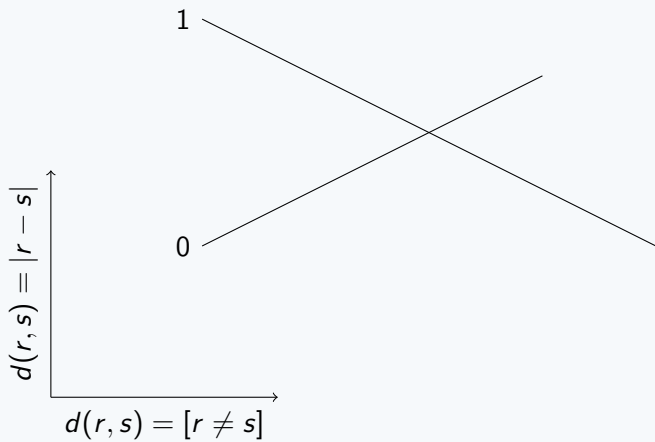
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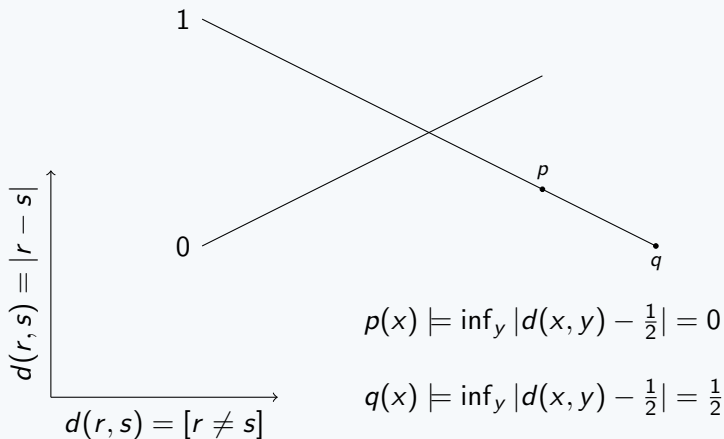
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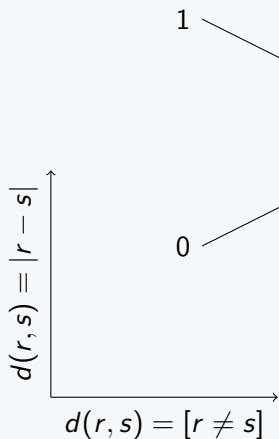
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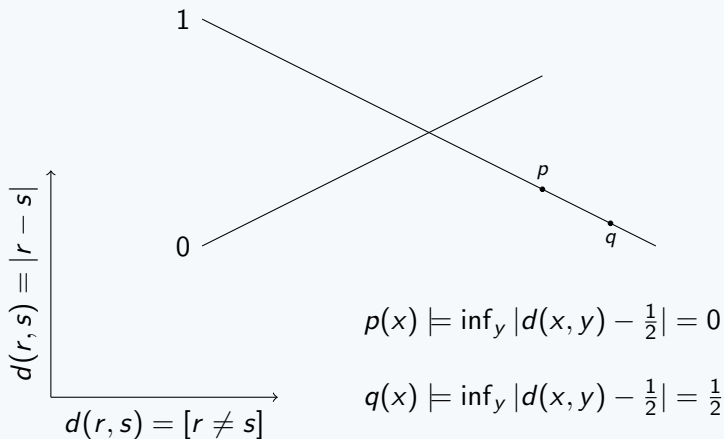


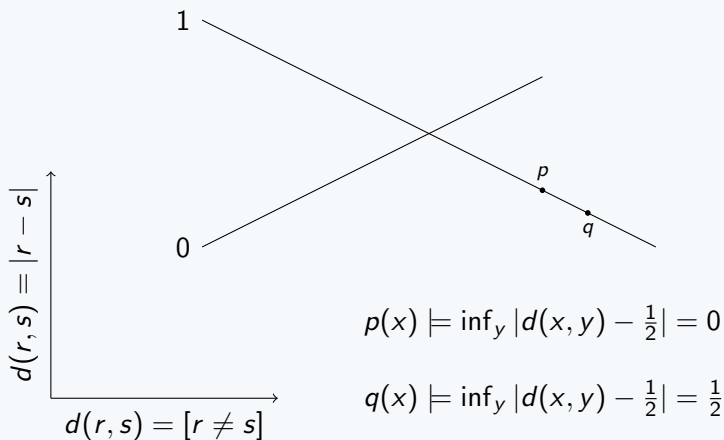


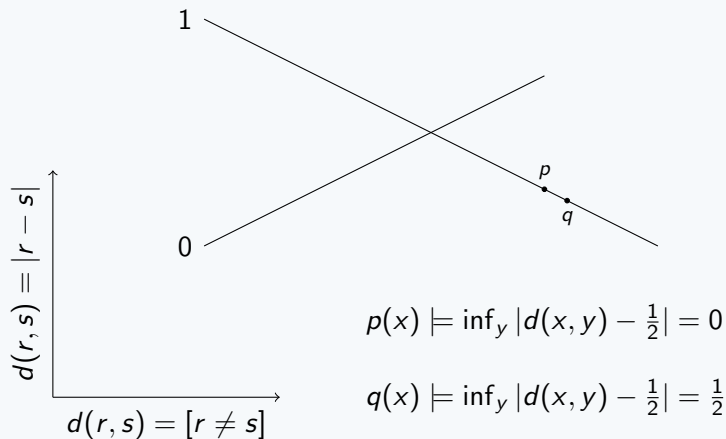


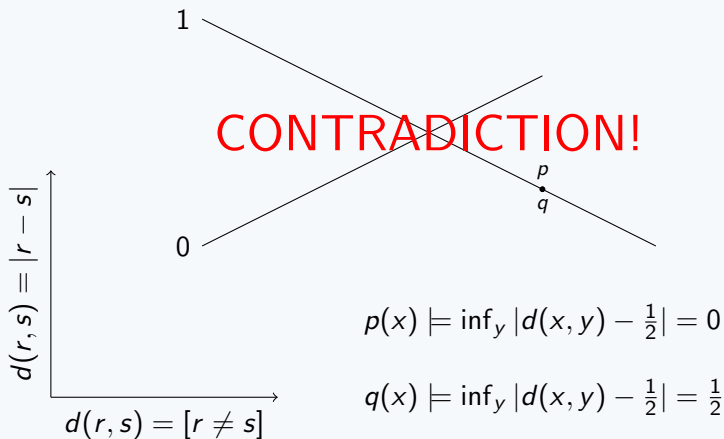
$$p(x) \models \inf_y |d(x, y) - \frac{1}{2}| = 0$$

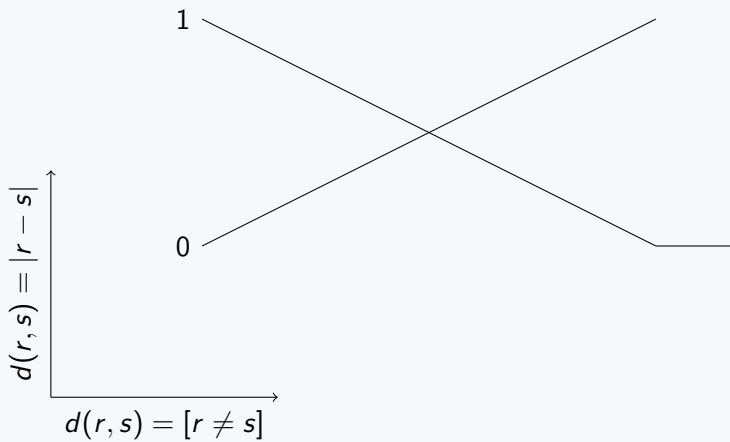
$$q(x) \models \inf_y |d(x, y) - \frac{1}{2}| = \frac{1}{2}$$

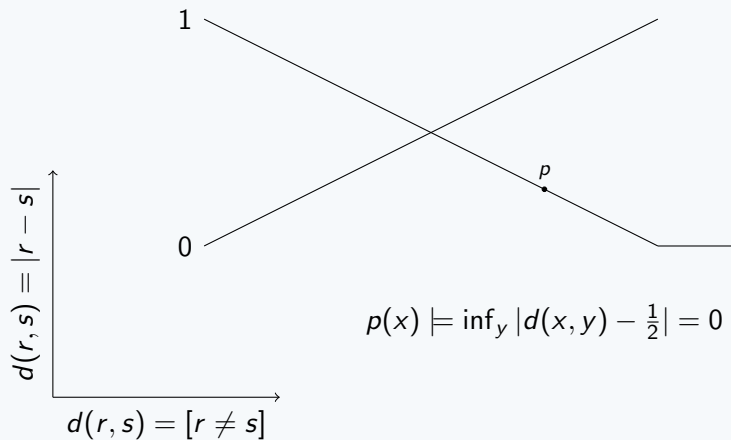


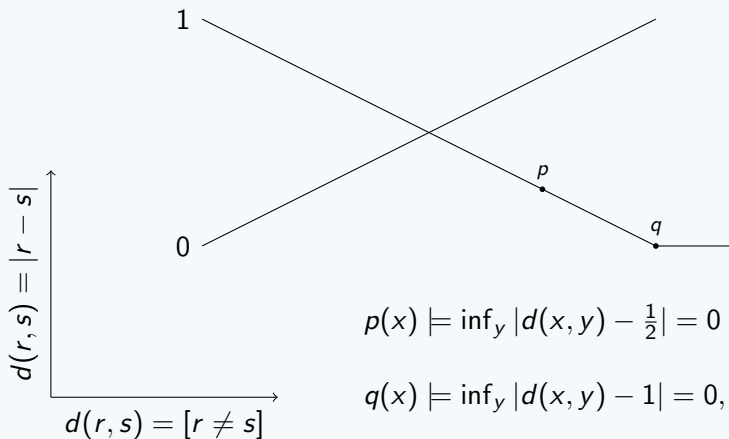


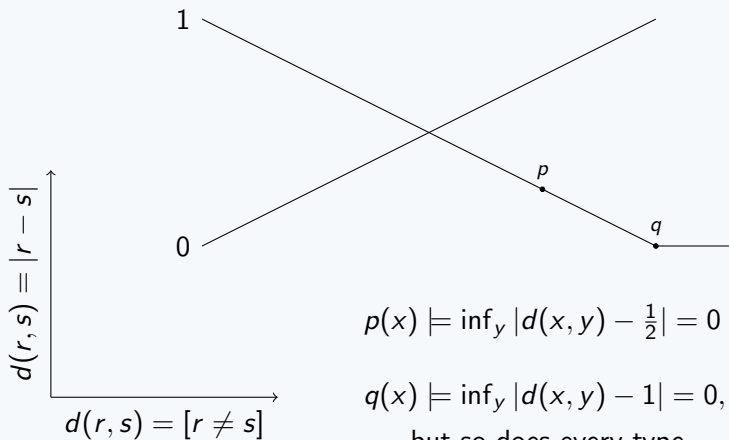








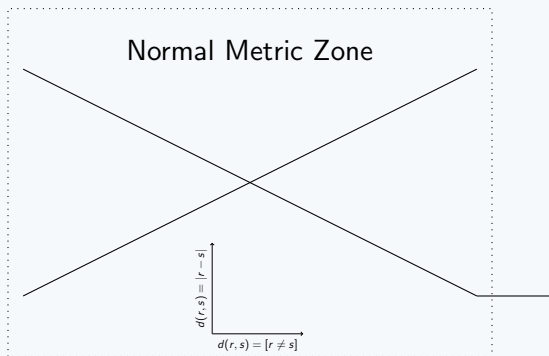




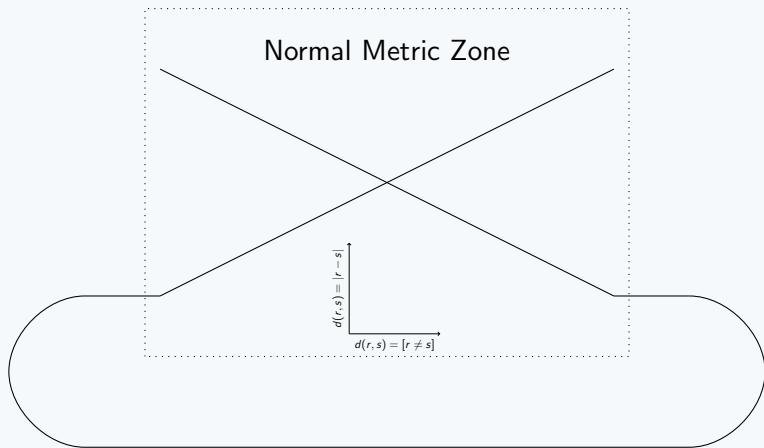
$$p(x) \models \inf_y |d(x, y) - \frac{1}{2}| = 0$$

$$q(x) \models \inf_y |d(x, y) - 1| = 0,$$

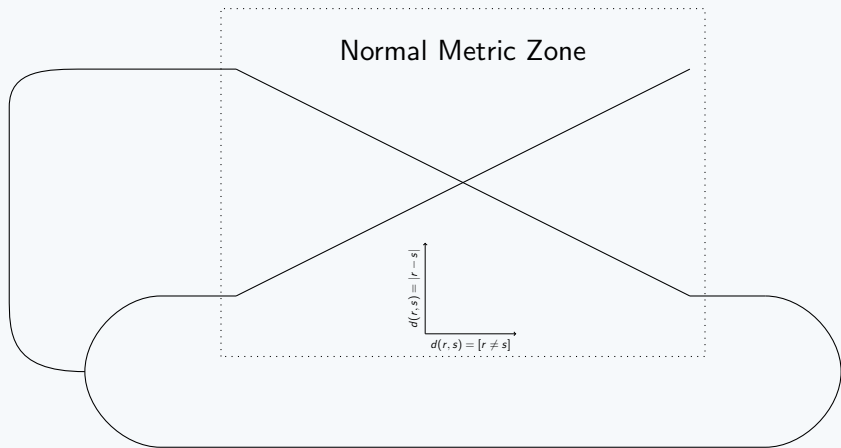
but so does every type.



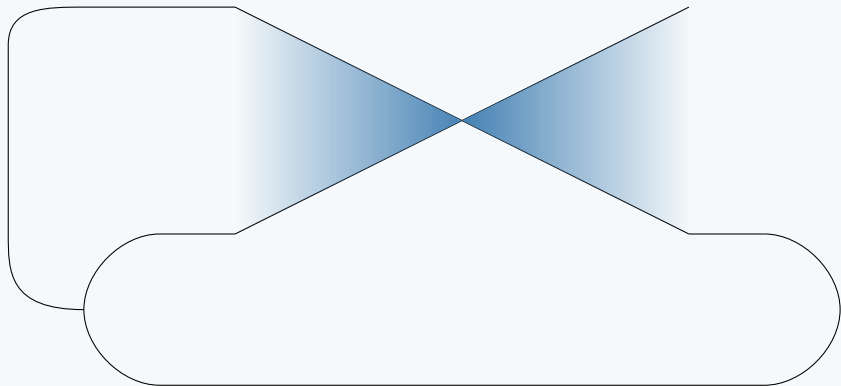
All Distances Are 0 or 1 Zone

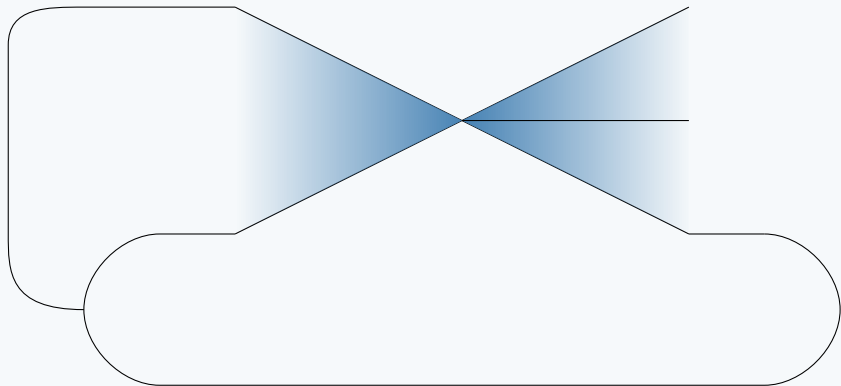


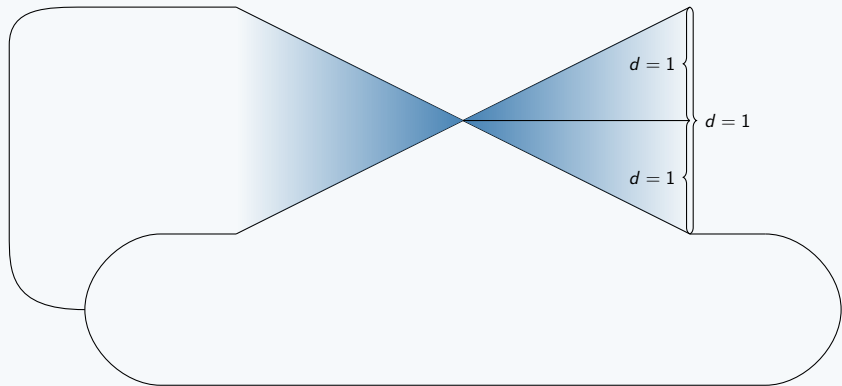
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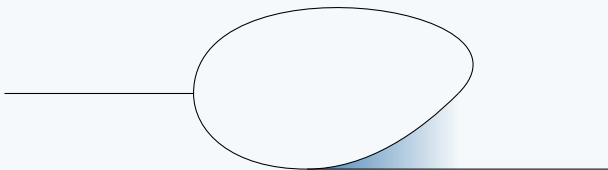
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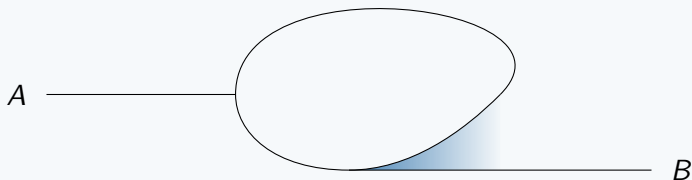


A diode



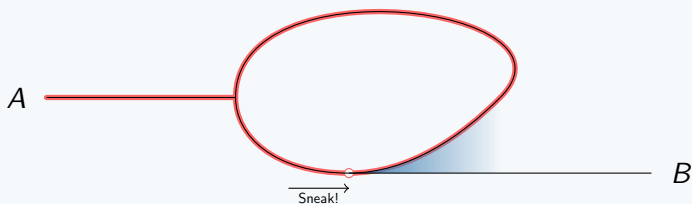
A diode

| Open Set U



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$S_1(T) \setminus U$ is not definable. **X**

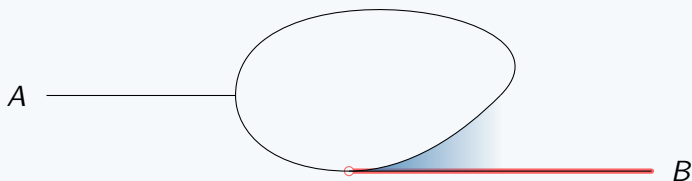
A diode

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$S_1(T) \setminus U$ is definable. ✓

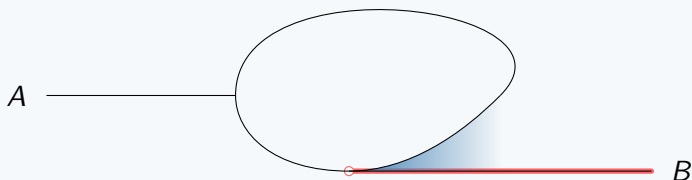
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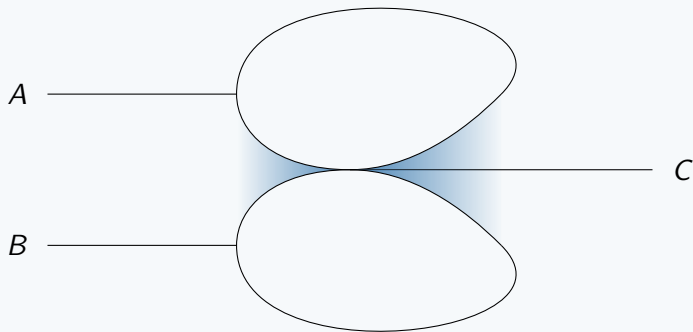
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Roughly: $S_1(T) \setminus U$ is definable iff $A \in U \rightarrow B \in U$.

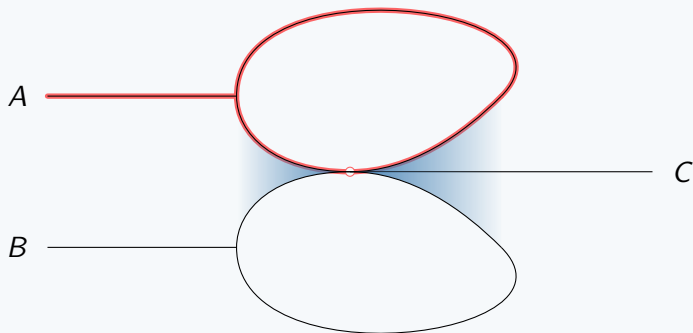
An AND gate

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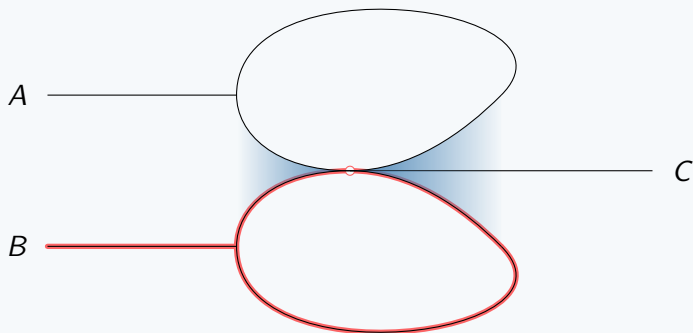
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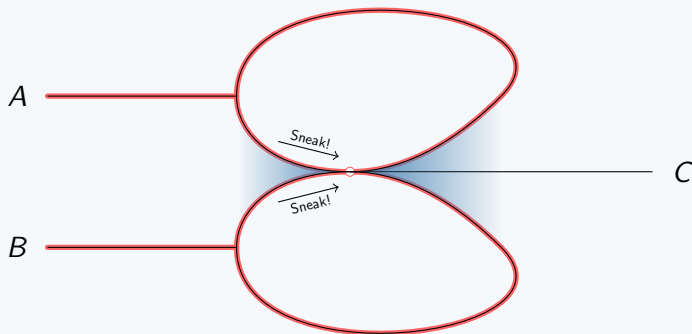
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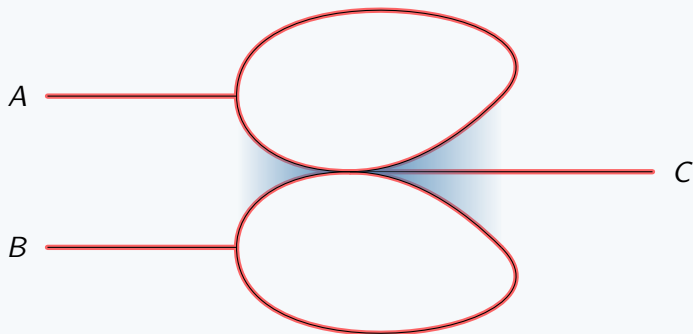
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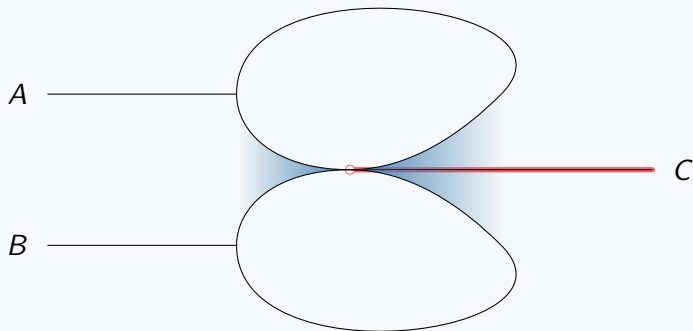
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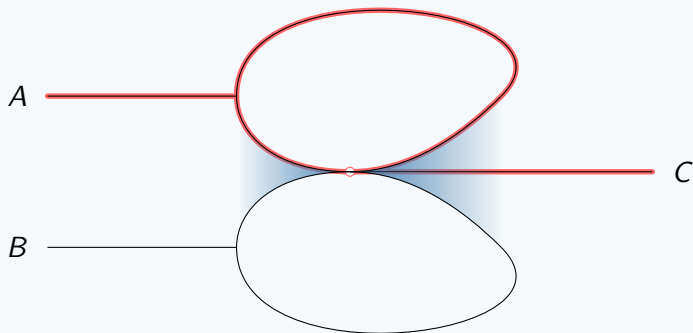
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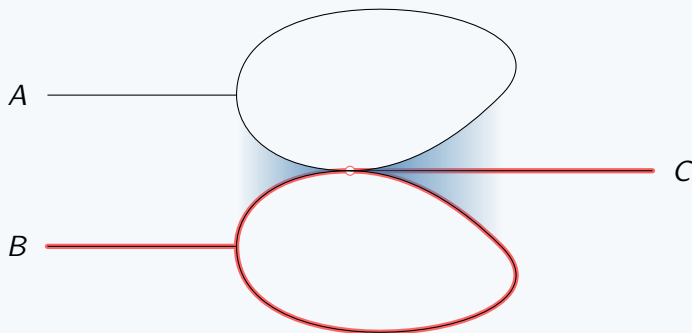
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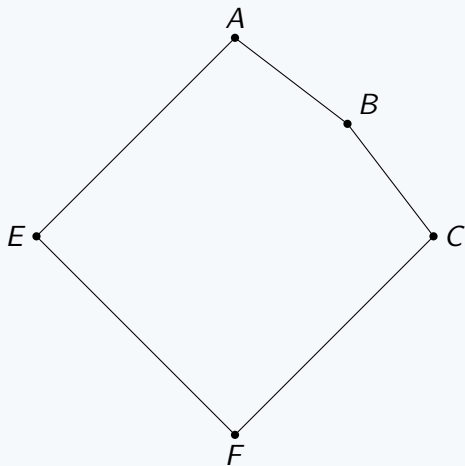
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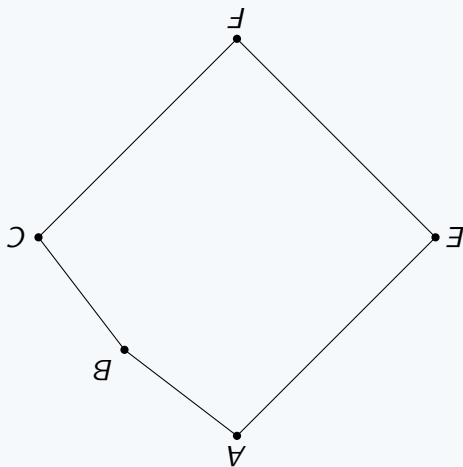
The construction

Take your favorite finite lattice with more than one element



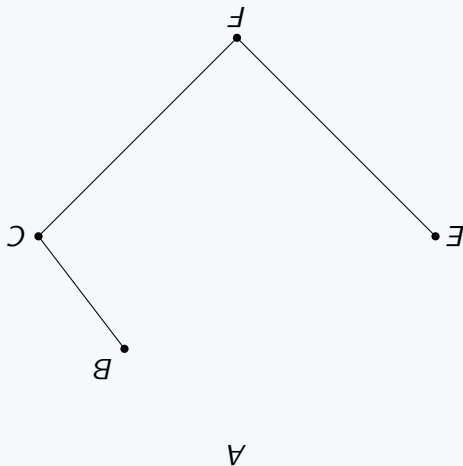
The construction

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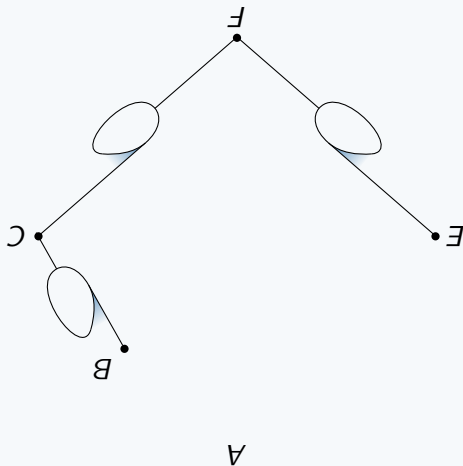
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pair with $x \leq y$

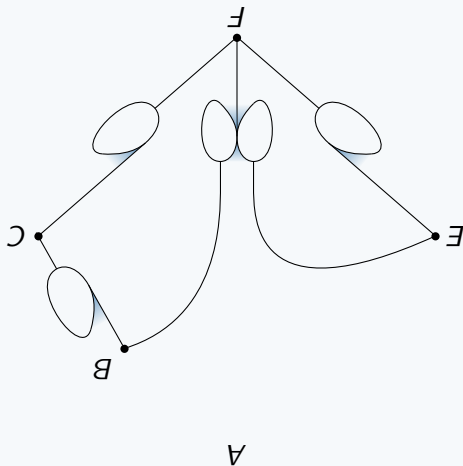


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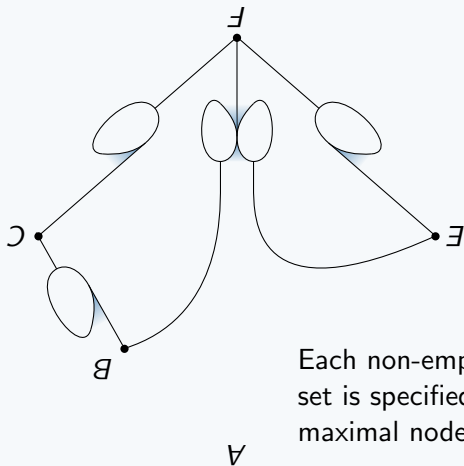


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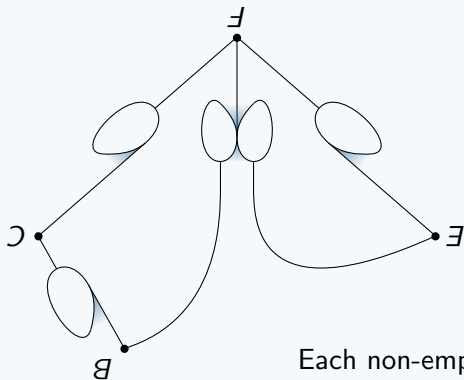
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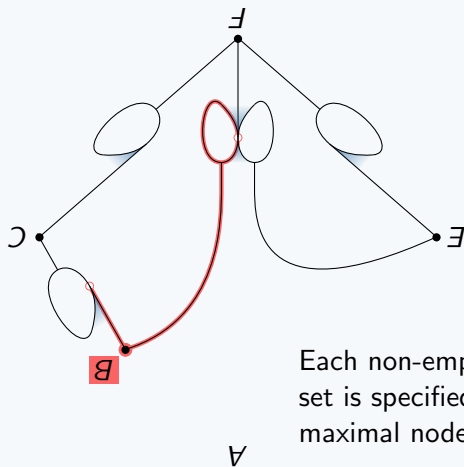
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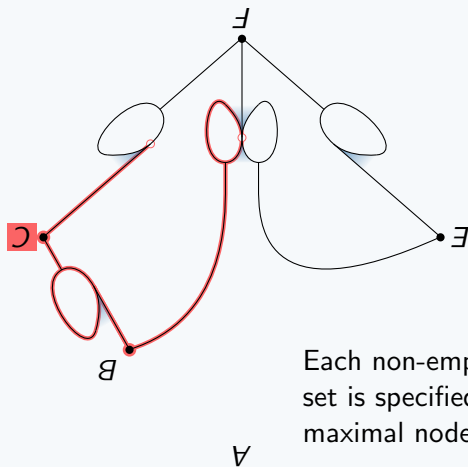
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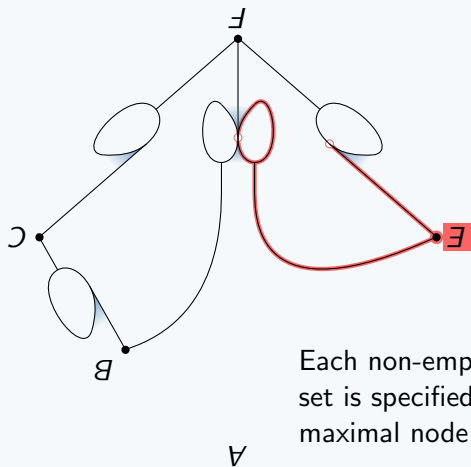
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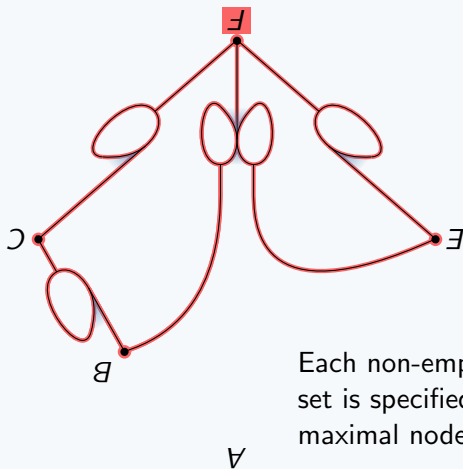
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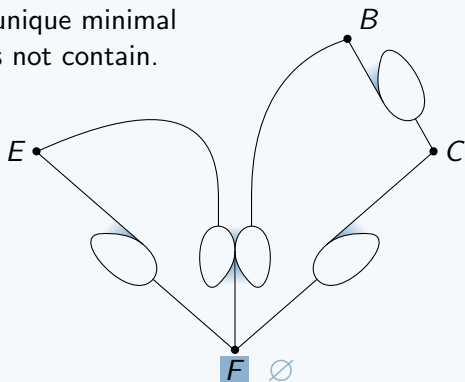
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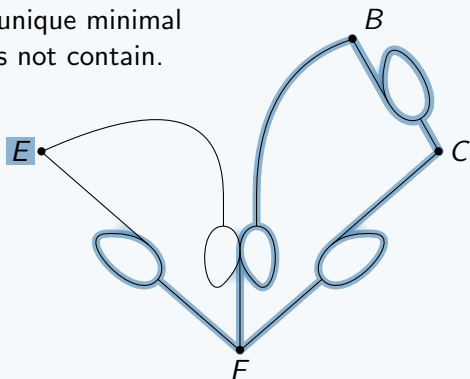


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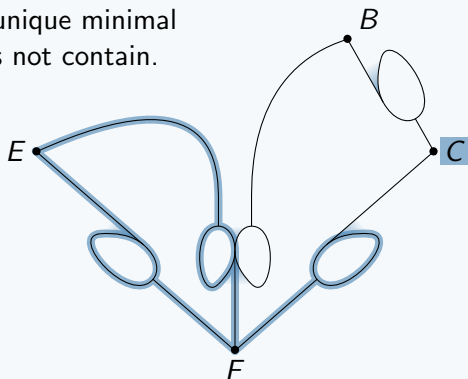


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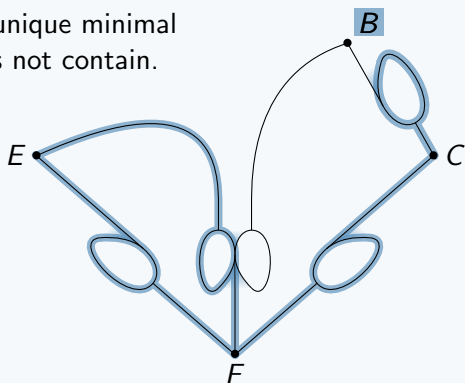


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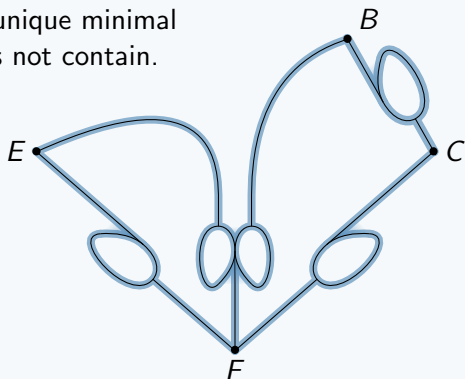


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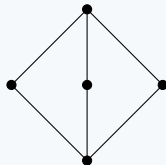
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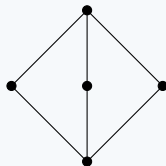
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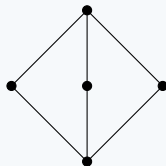
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Question

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Some infinite lattices

Which infinite semilattices can we have? I

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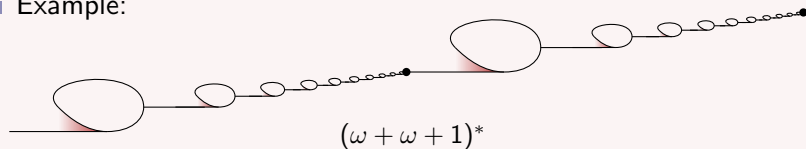
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Thank you